

Weaknesses in Defenses against Web-Borne Malware

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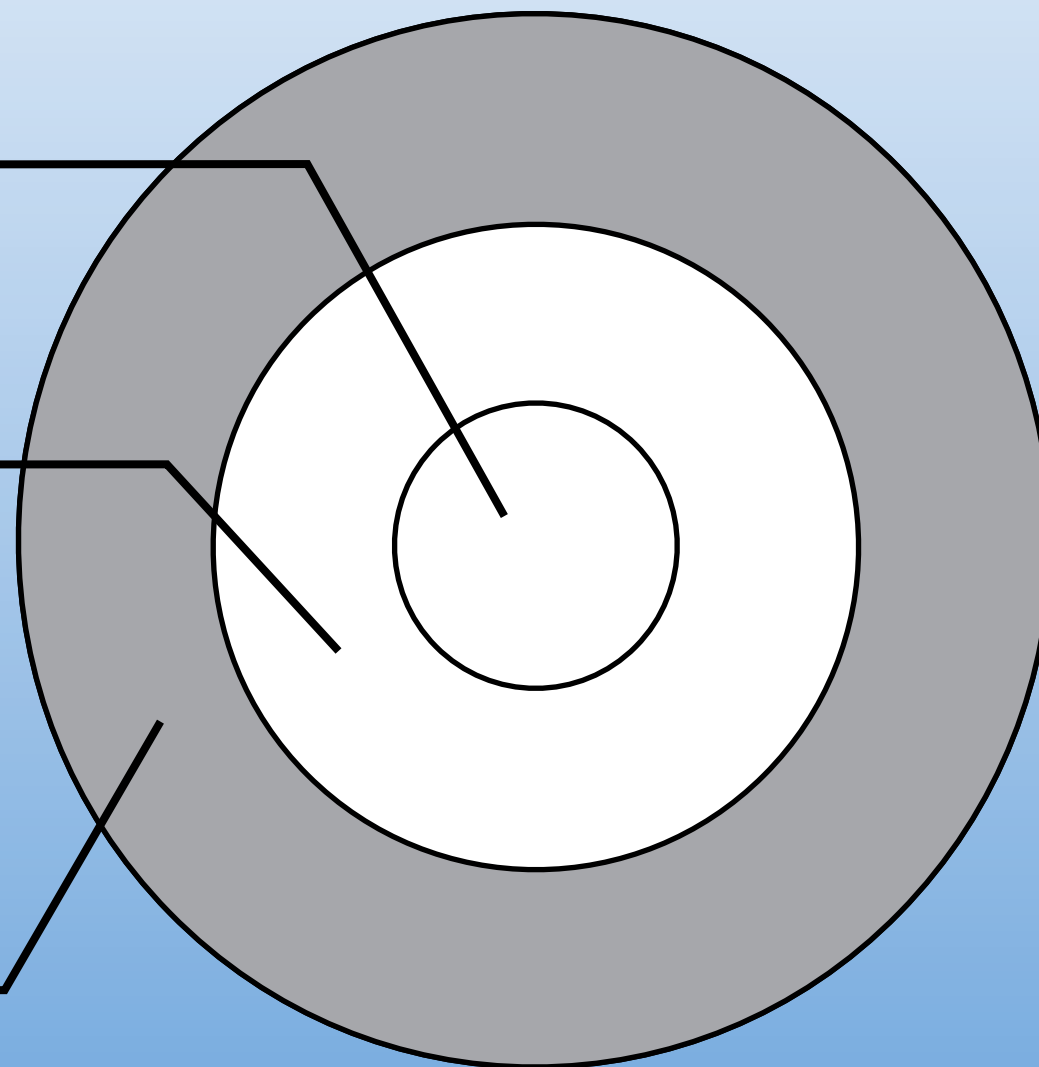
July, 2013

Web-based Malware

Payload

Infiltration

Protection



Motivation

Existing detectors are tied to today's malware protection techniques

- *Can we get around the detectors by breaking the assumptions they made?*

Outline

- Current techniques
- Thwarting detection/analysis
- Experimental evaluation
- Discussion

Current Techniques

Static Analysis

Zozzle:

Curtsinger et al., 2011

Cujo:

Rieck et al., 2010

Assumption:

Malware only apply
code unpacking
obfuscation

Dynamic Analysis

JSAND:

Cova et al., 2010

Cujo:

Rieck et al., 2010

Assumption:

Malicious behavior
can be observed
under monitoring

Multi-Path Exploration

Rozzle:

Kolbitsch et al., 2012

Kudzu:

Saxena et al., 2010

Assumption:

Branching is done
by conditional
statements

Thwarting Analysis

Defense

Static analysis

simple deobfuscation

hiding

malicious logic

Dynamic analysis

single execution path

hiding

malicious behavior

Multi-path exploration

explore conditional
statements

hiding

branch logic

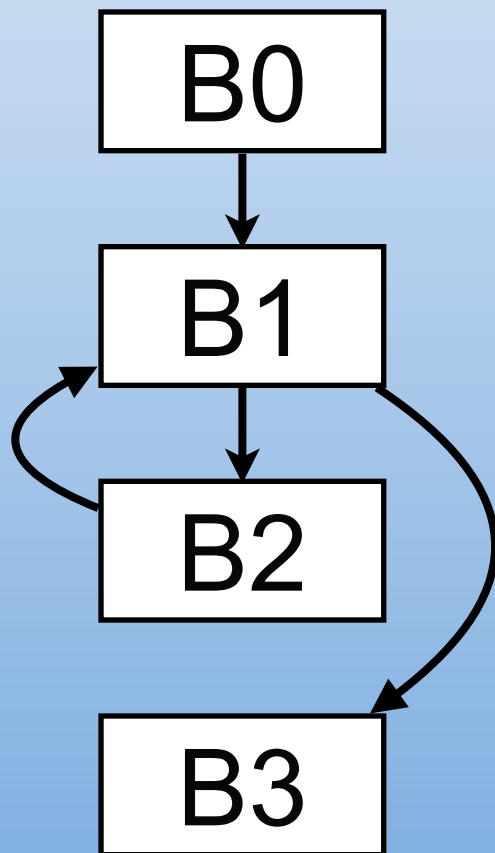
Exploitation

Emulation-Based
Obfuscation

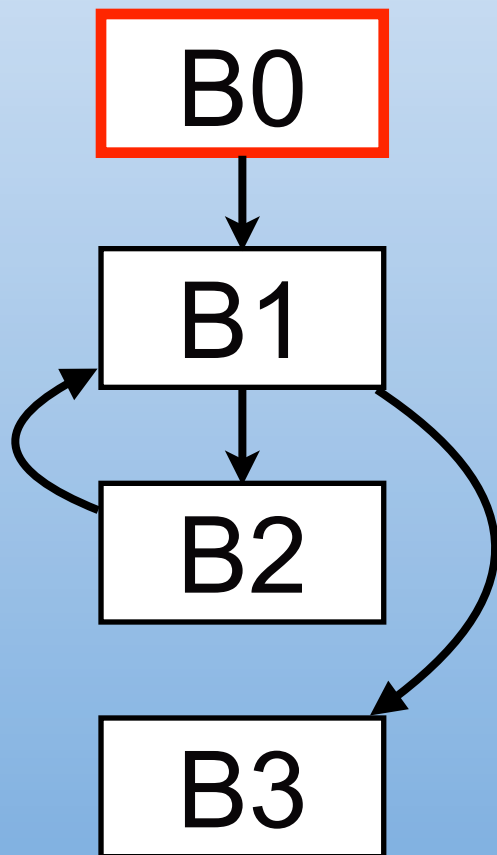
Anti-Analysis
Defense

Implicit
Conditionals

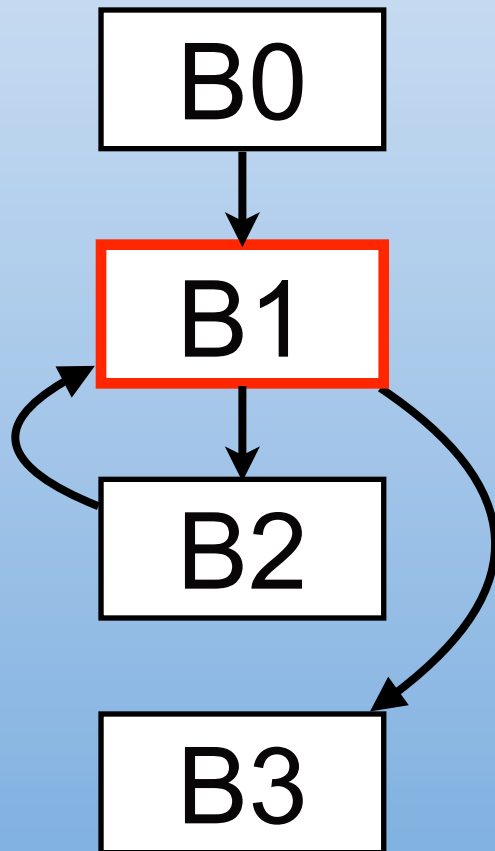
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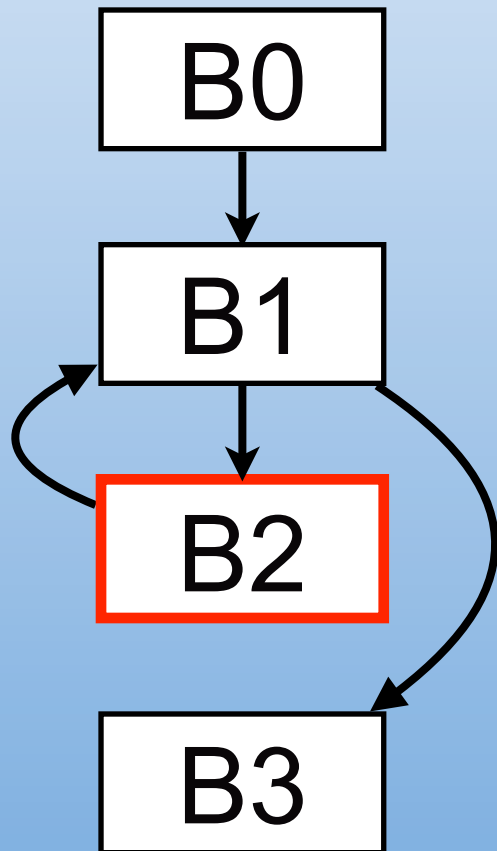
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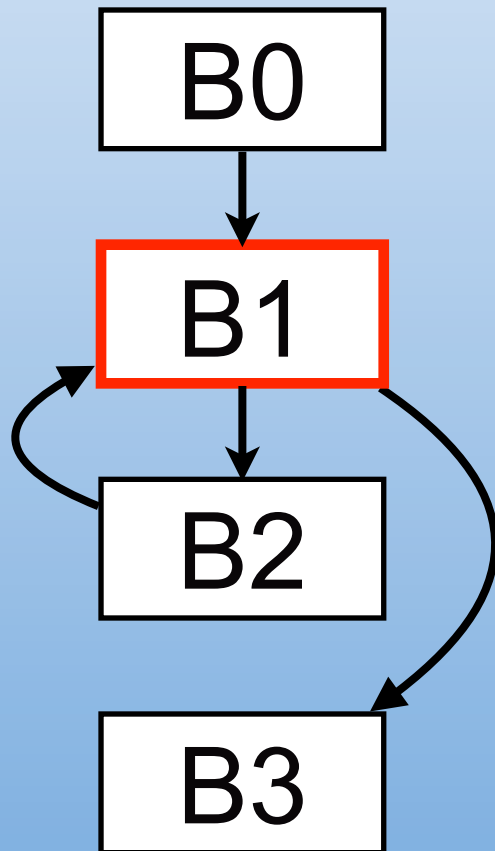
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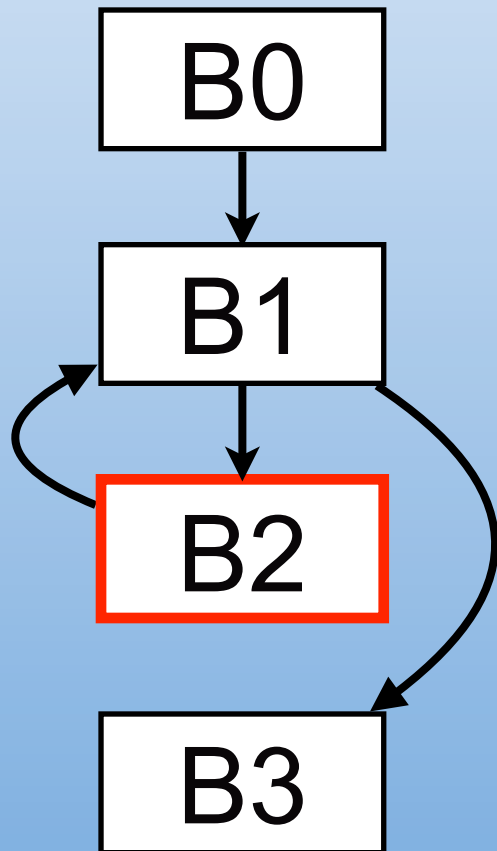
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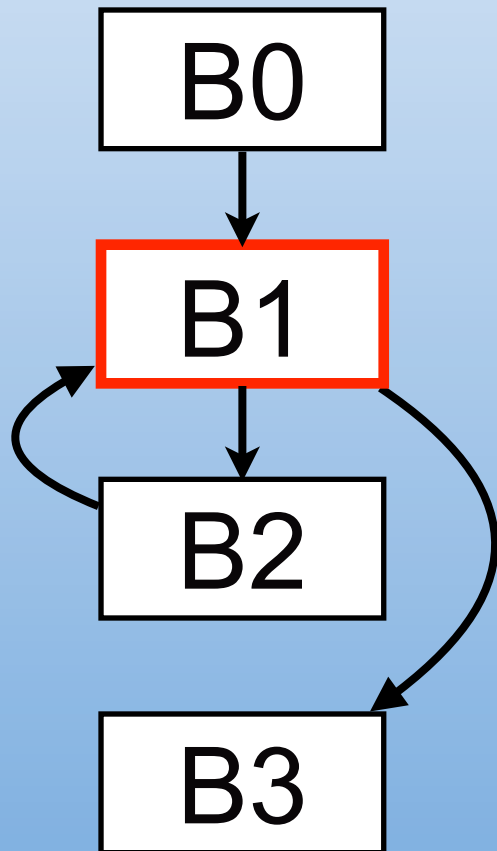
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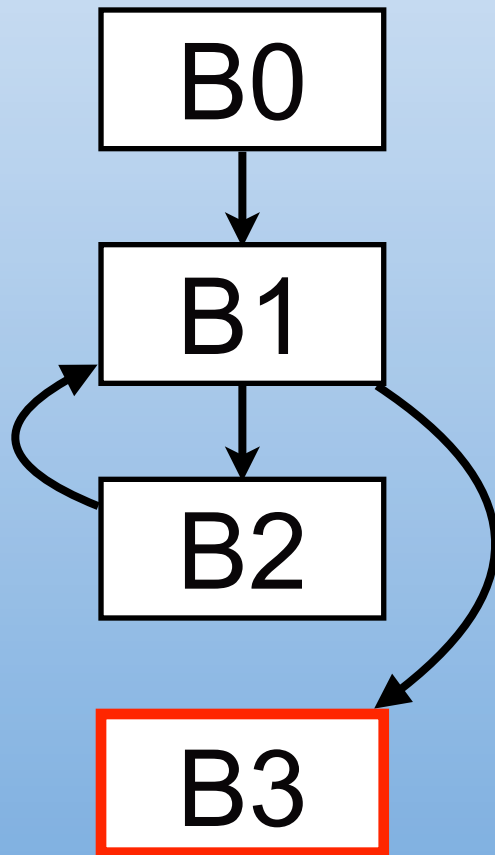
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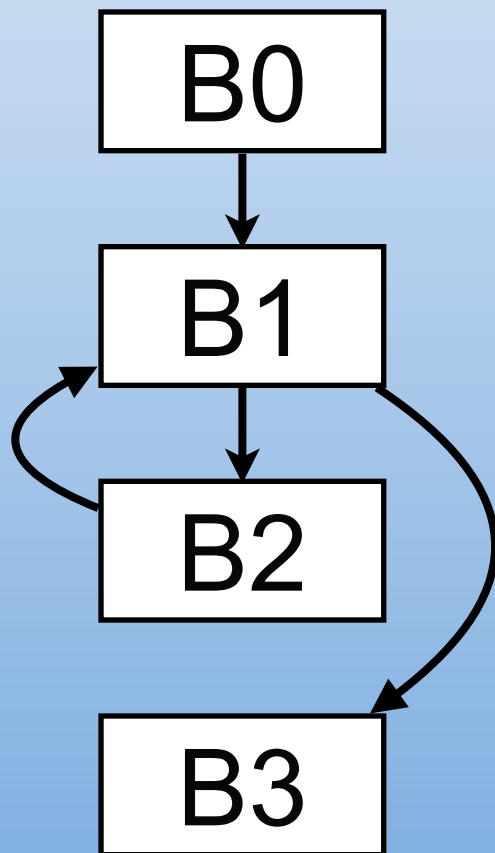
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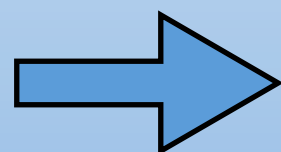
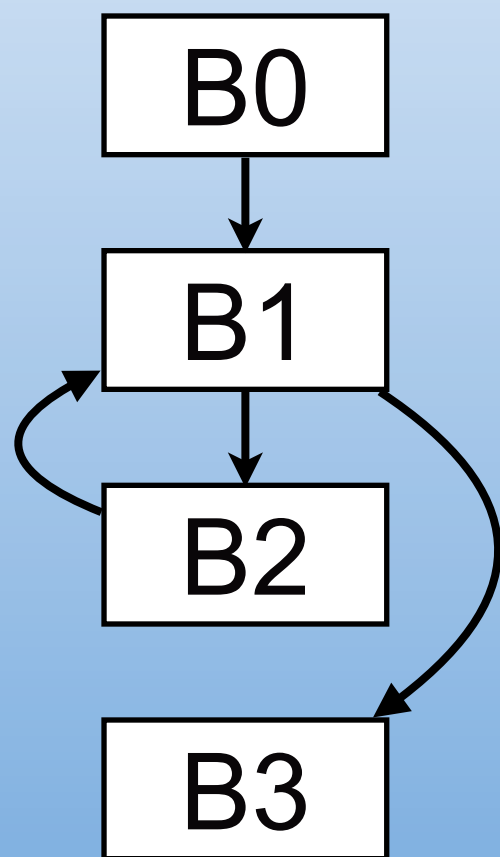
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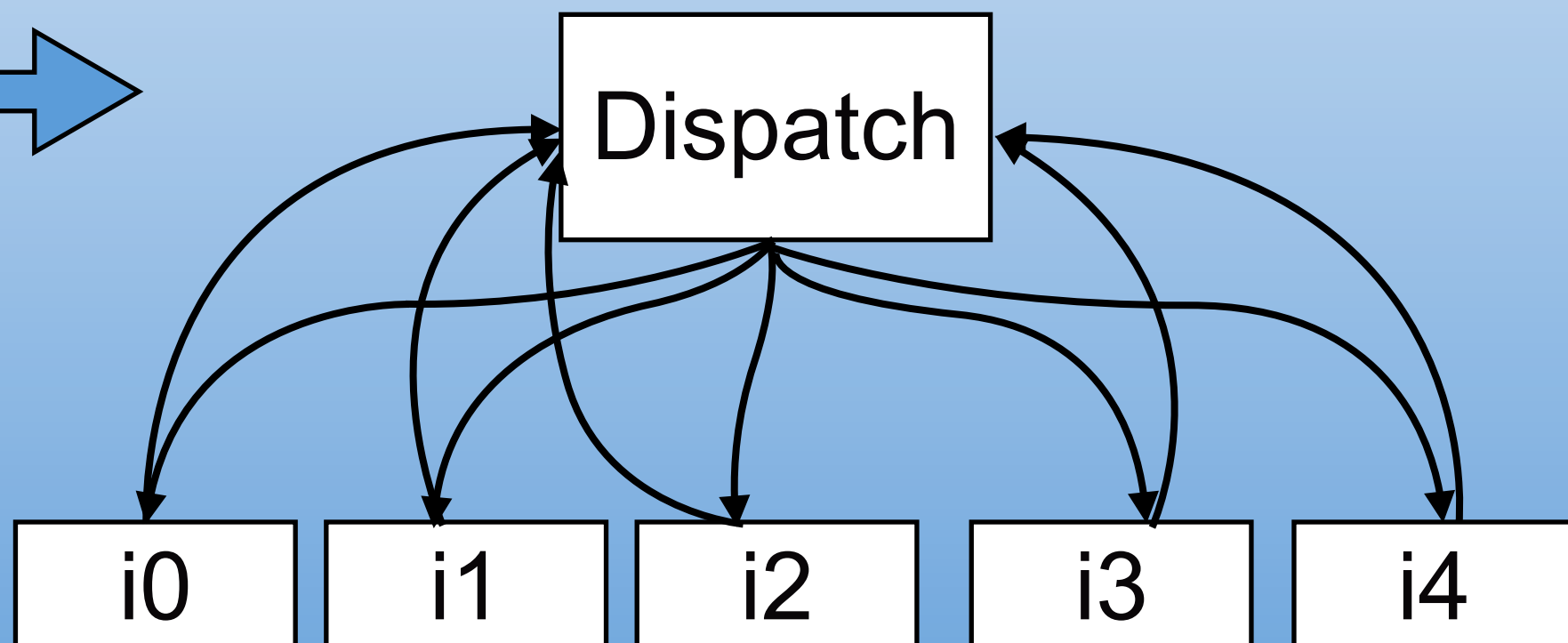
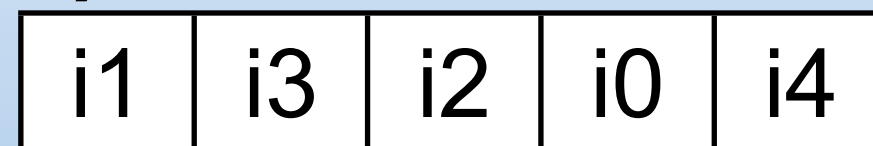
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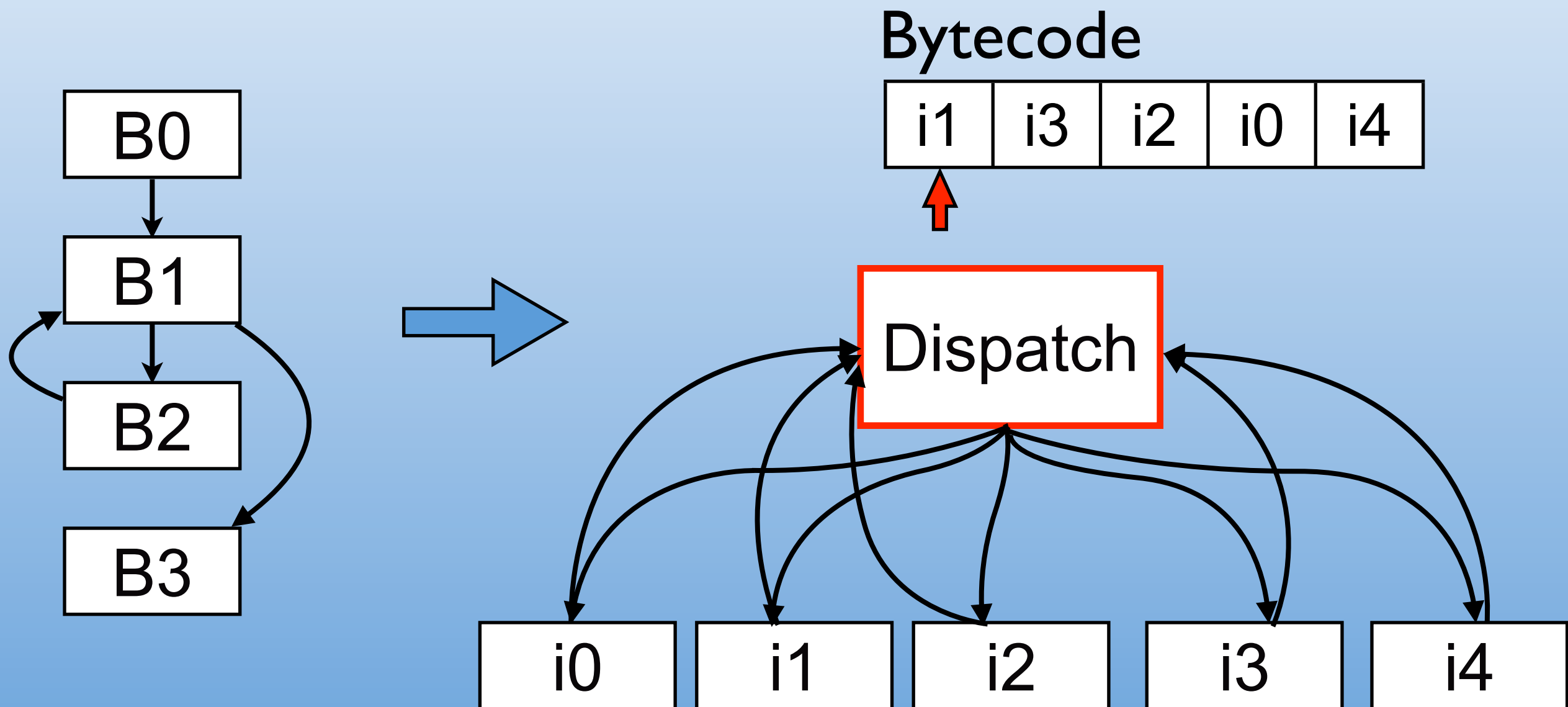
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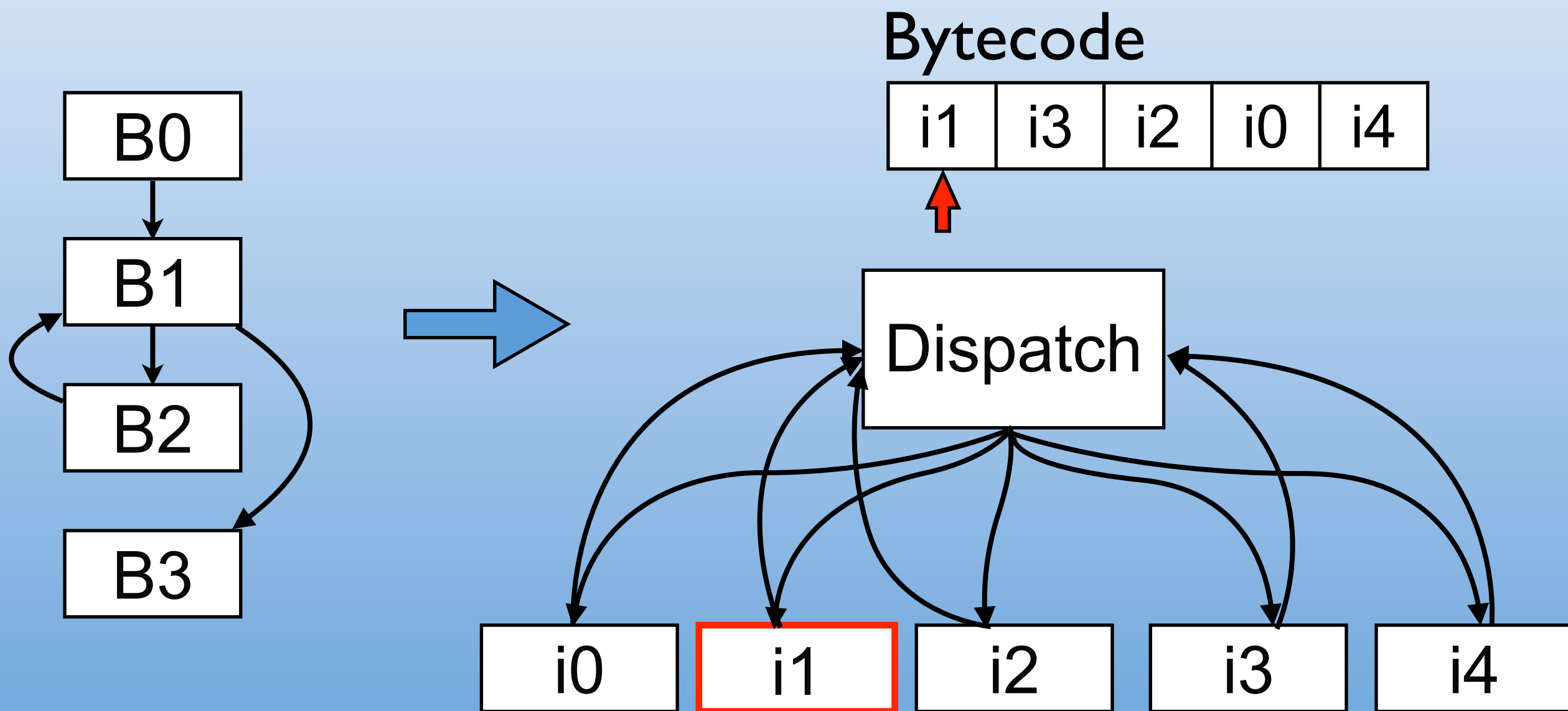
Bytecode



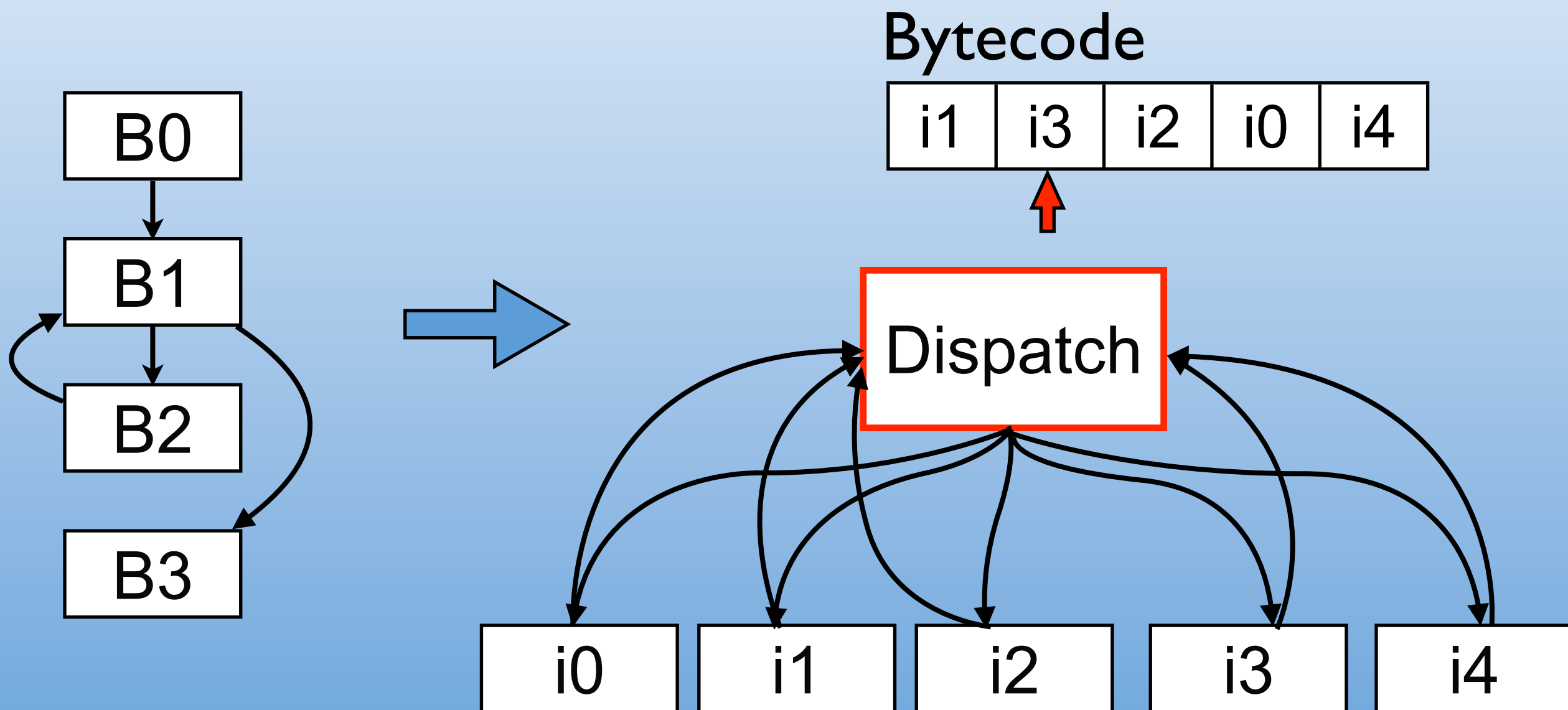
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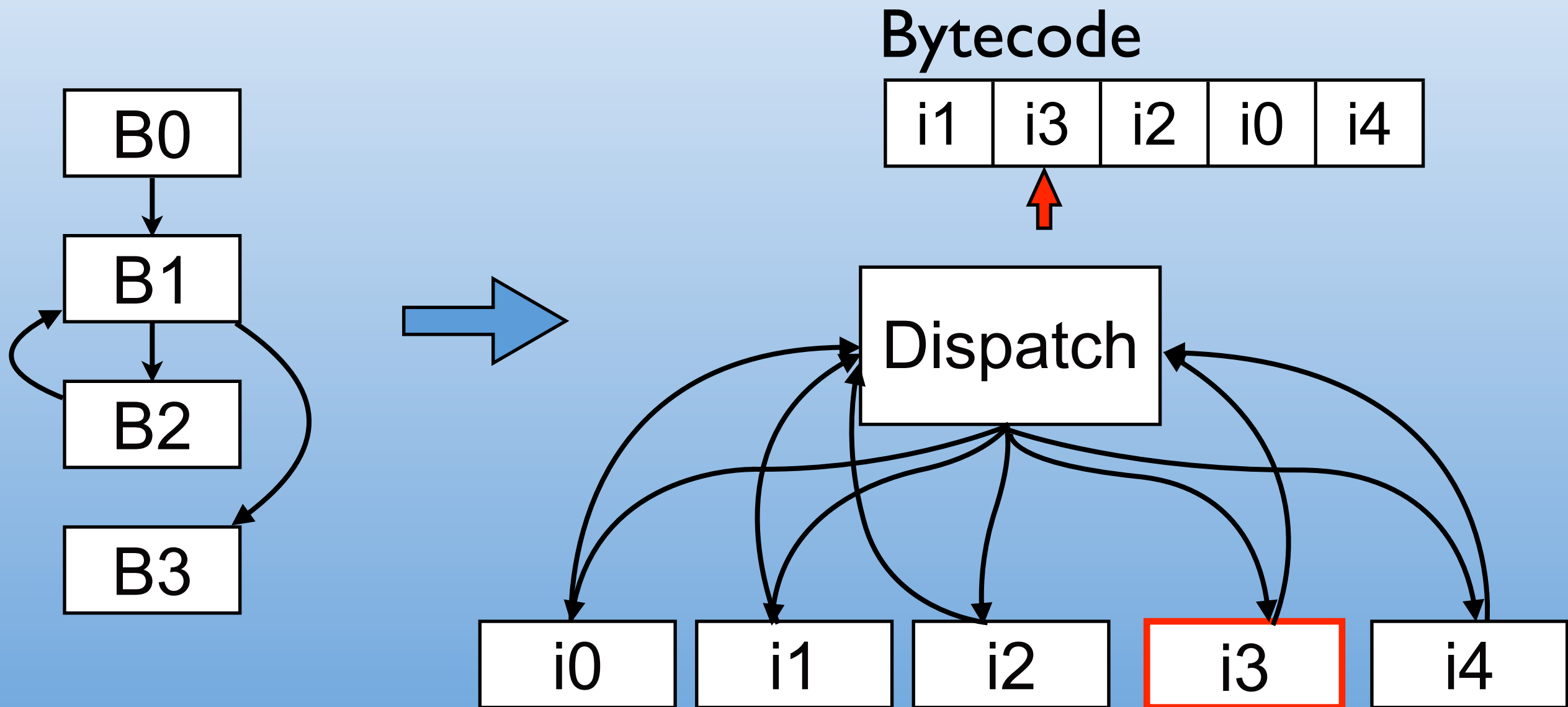
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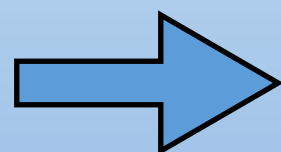
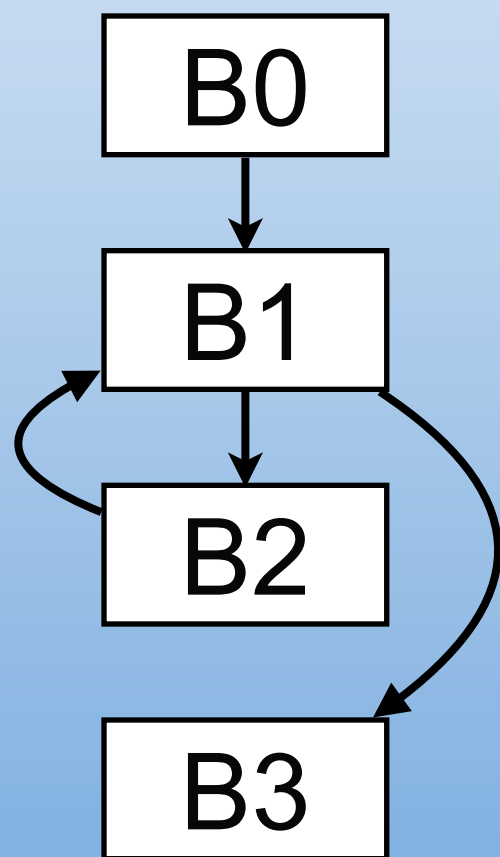
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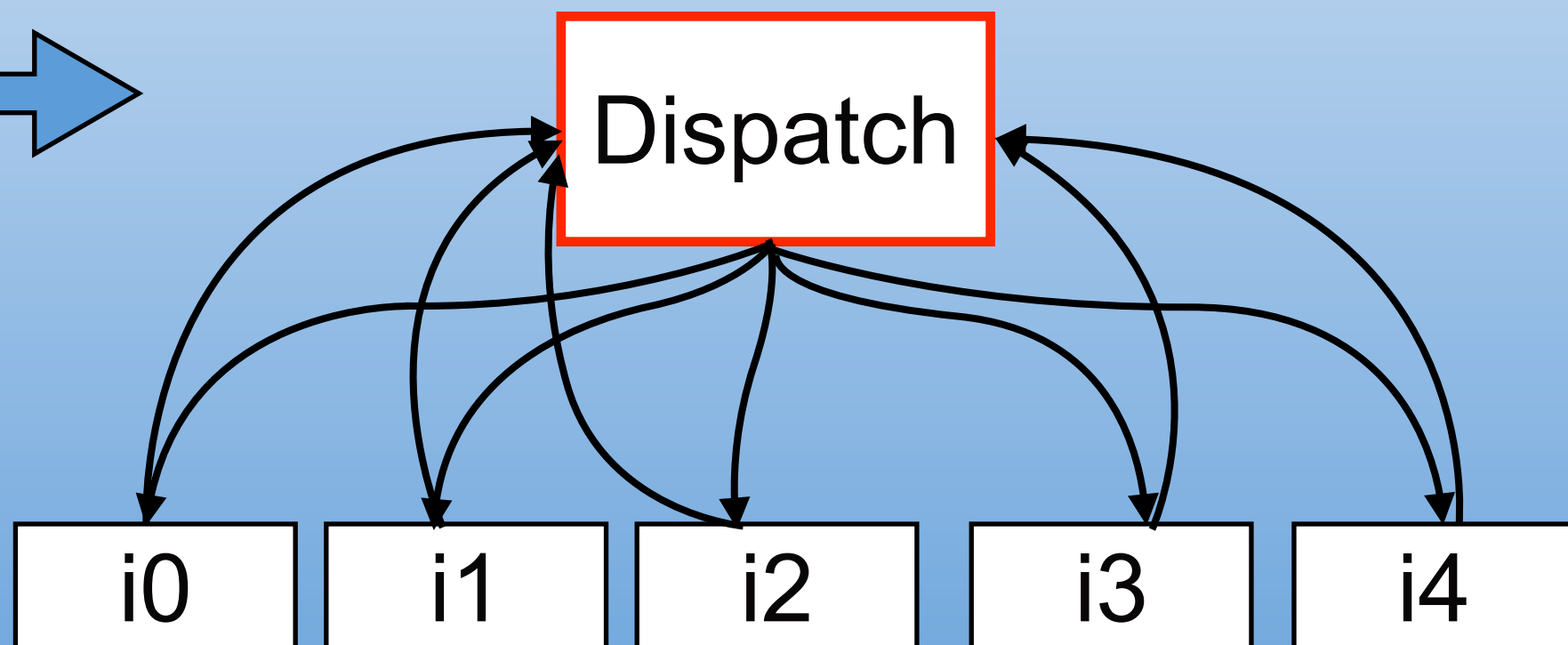
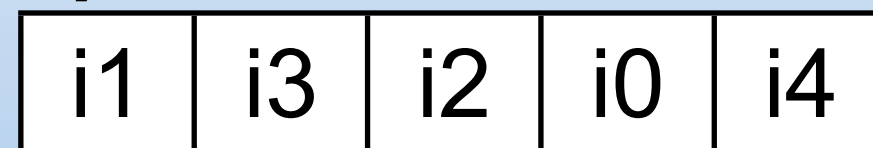
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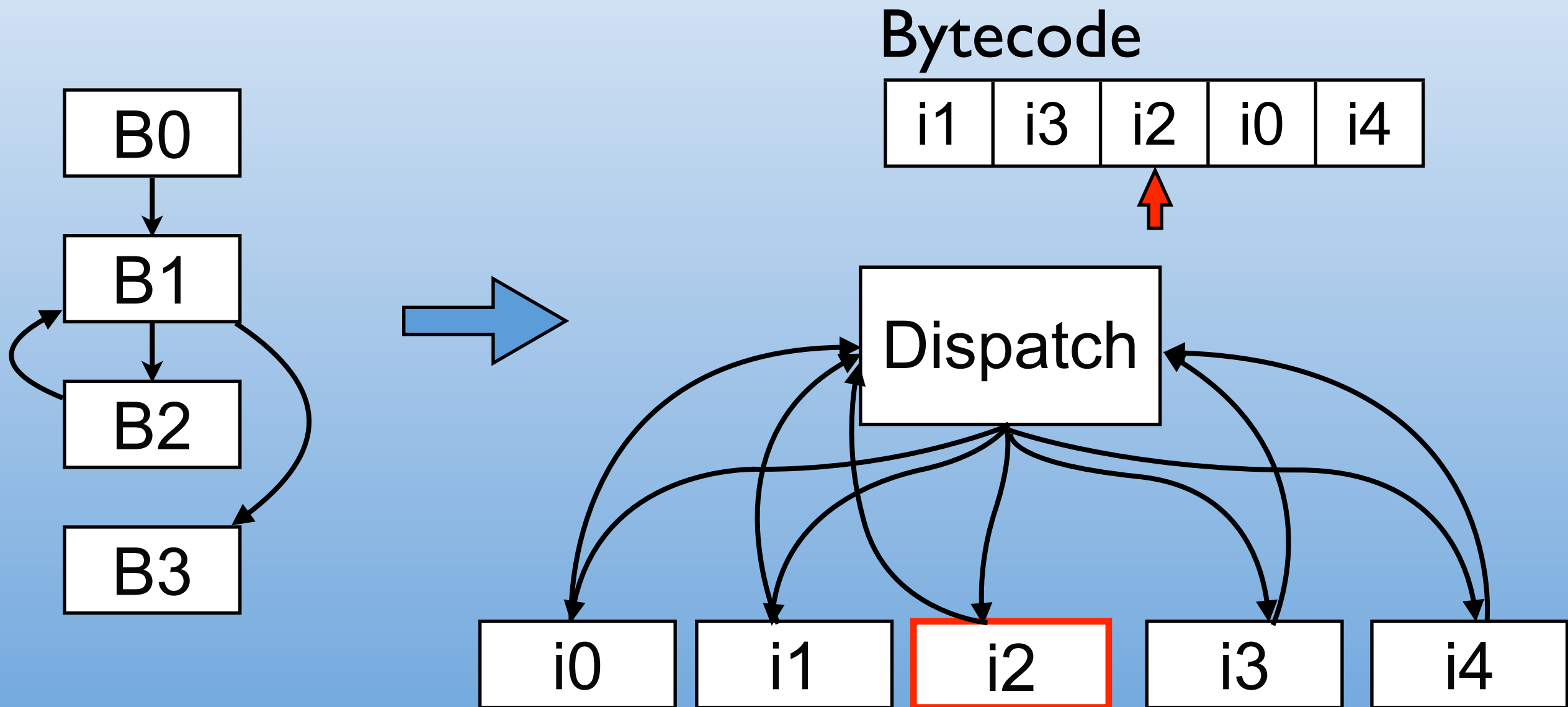
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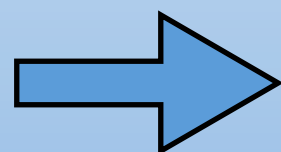
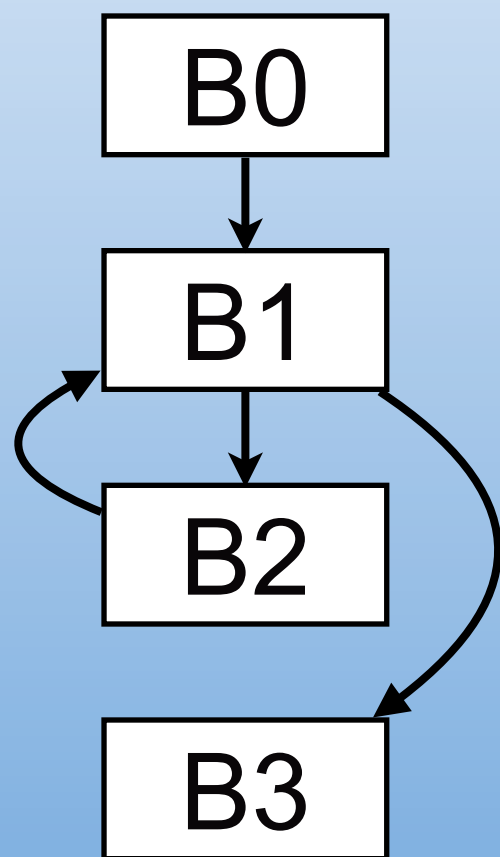
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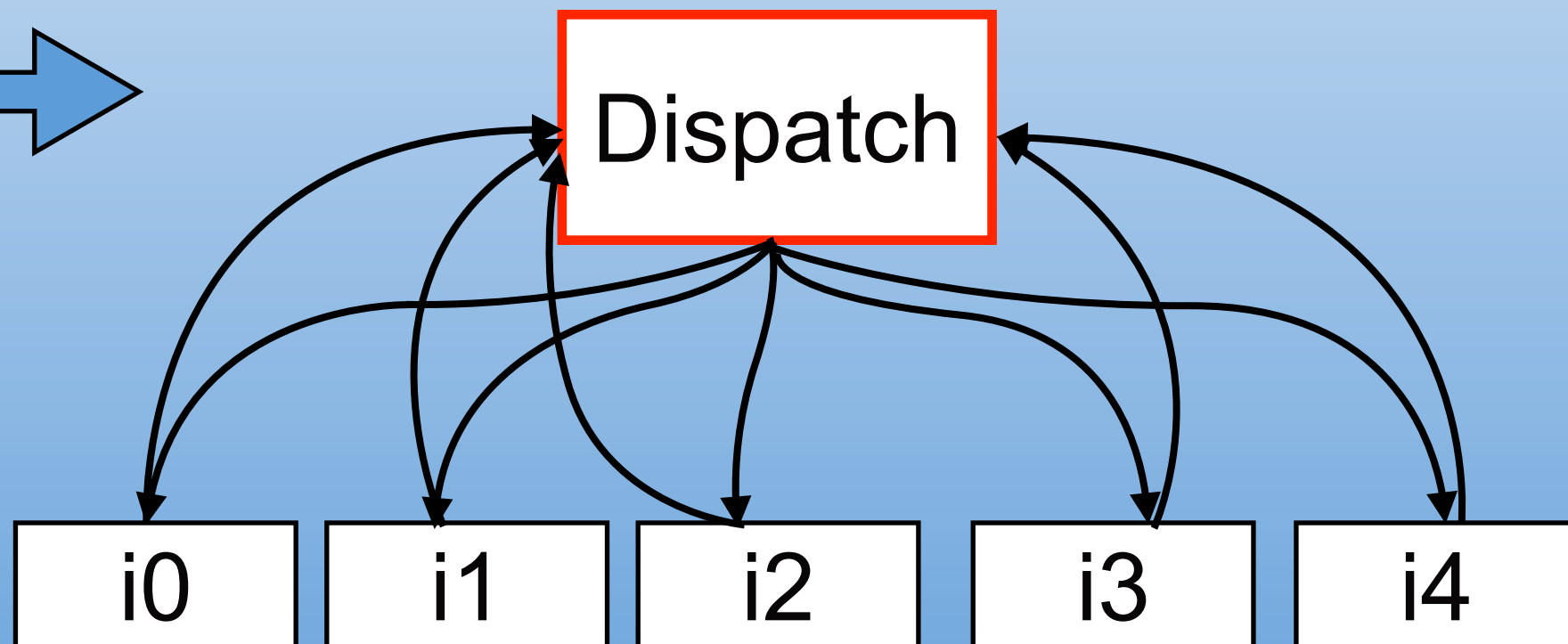
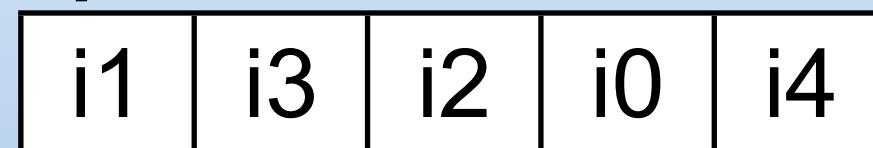
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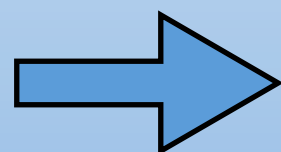
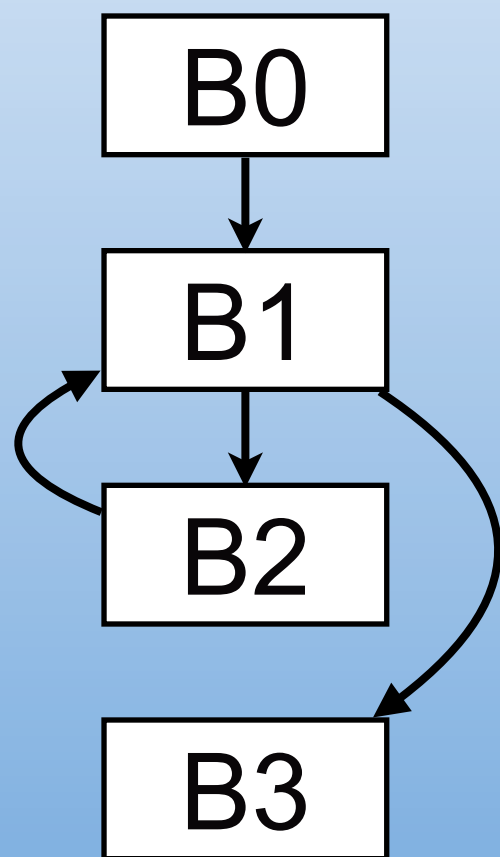
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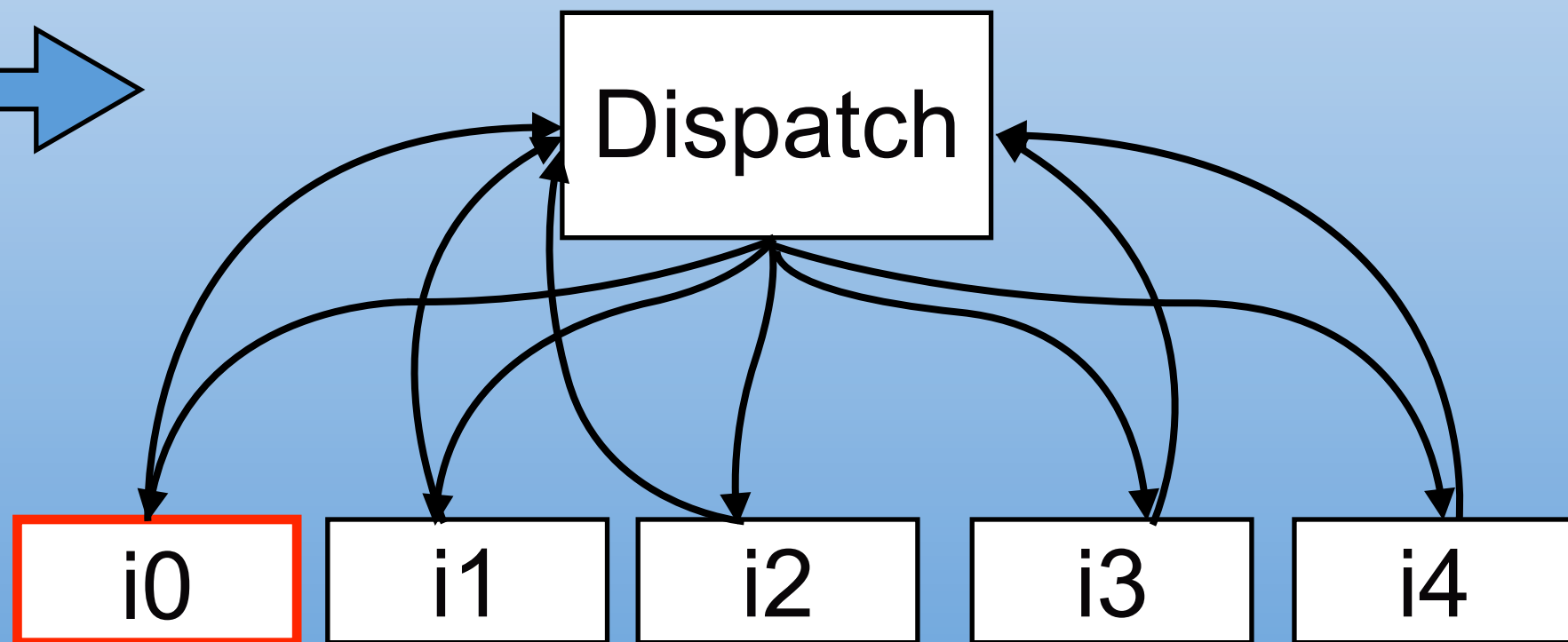
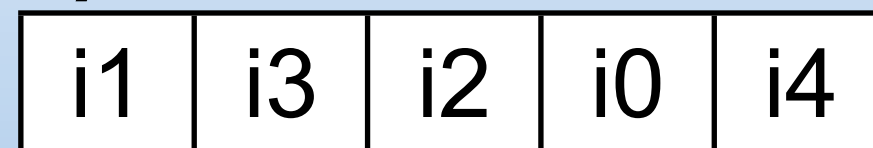
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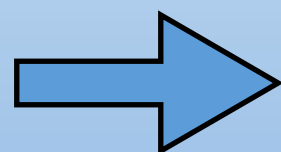
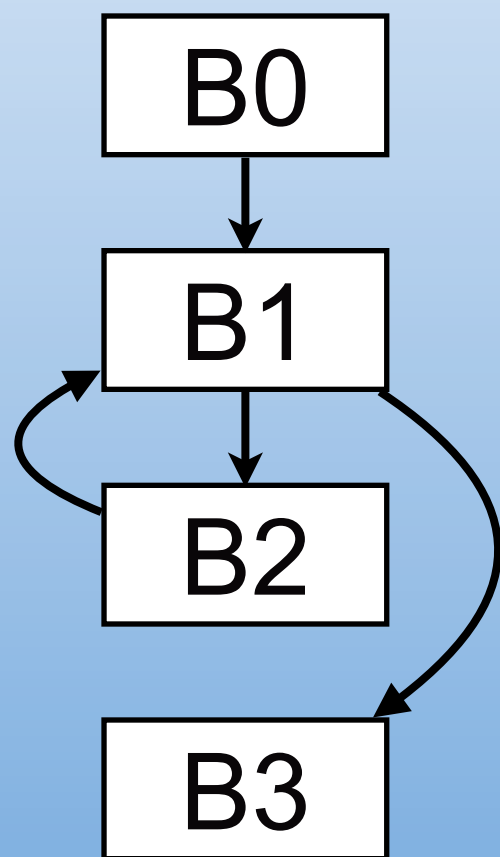
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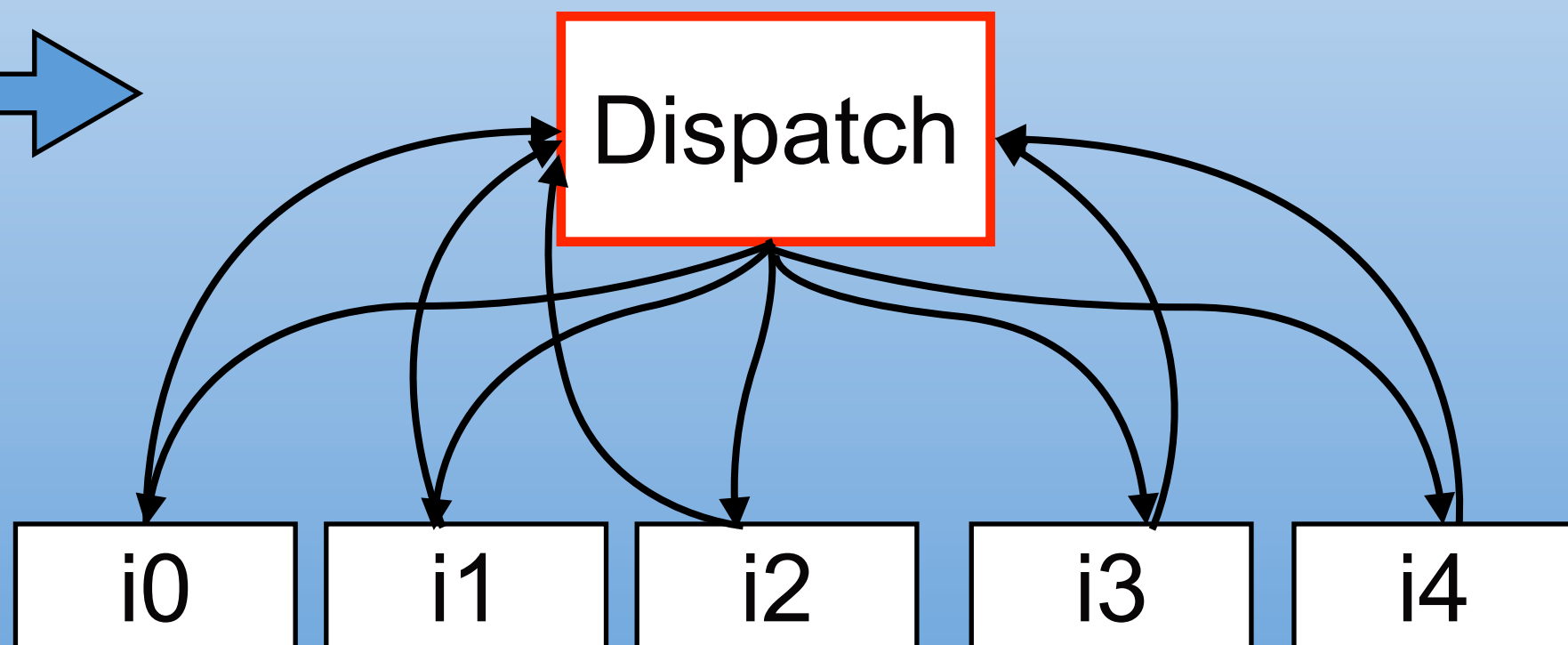
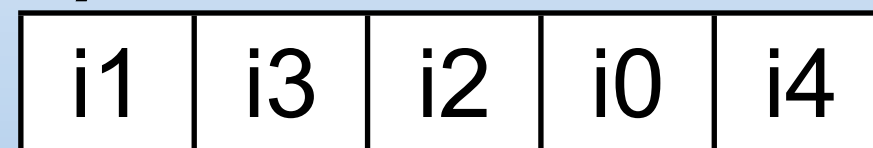
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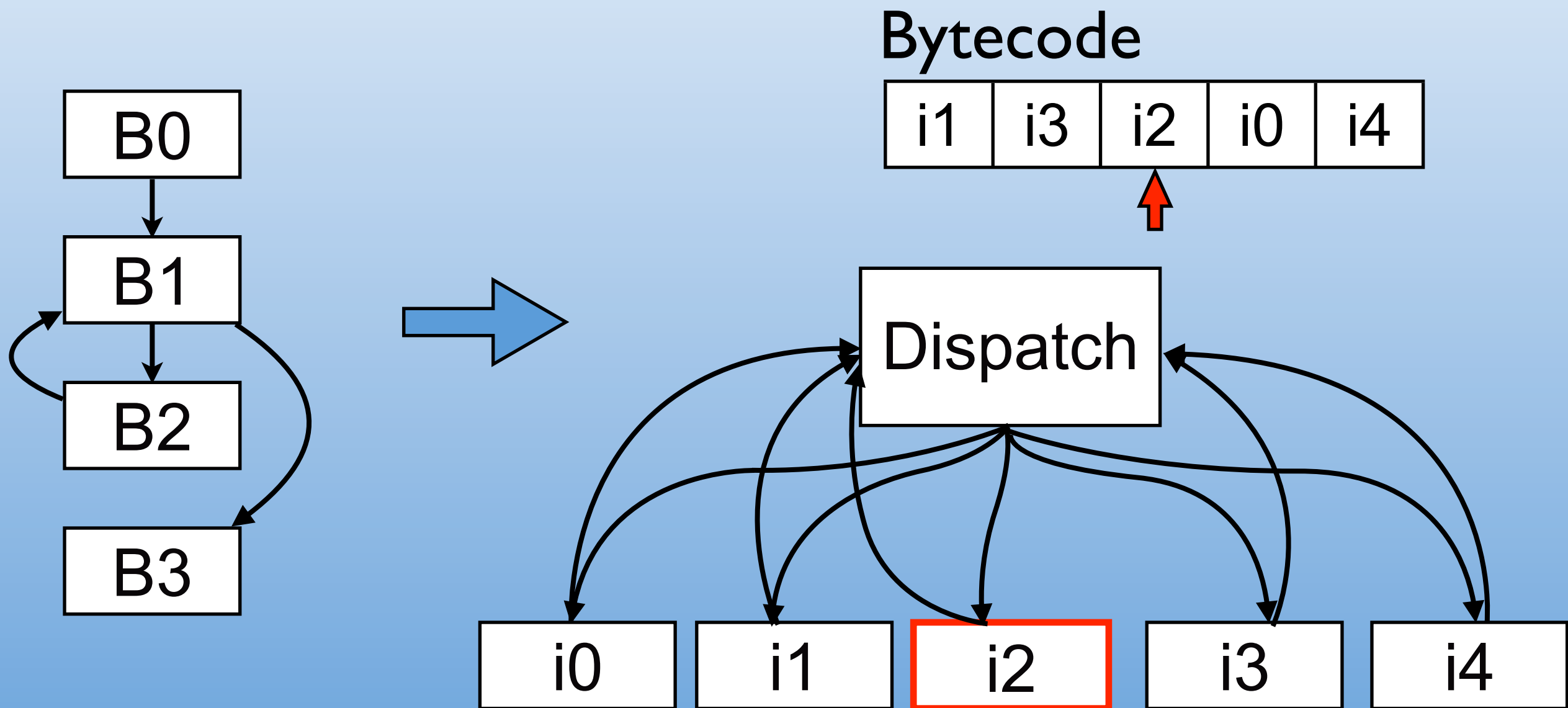
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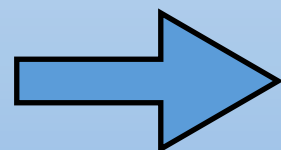
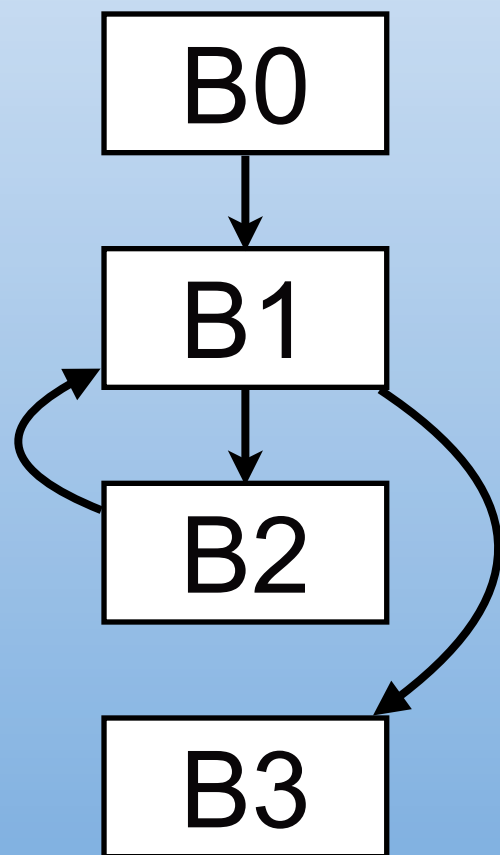
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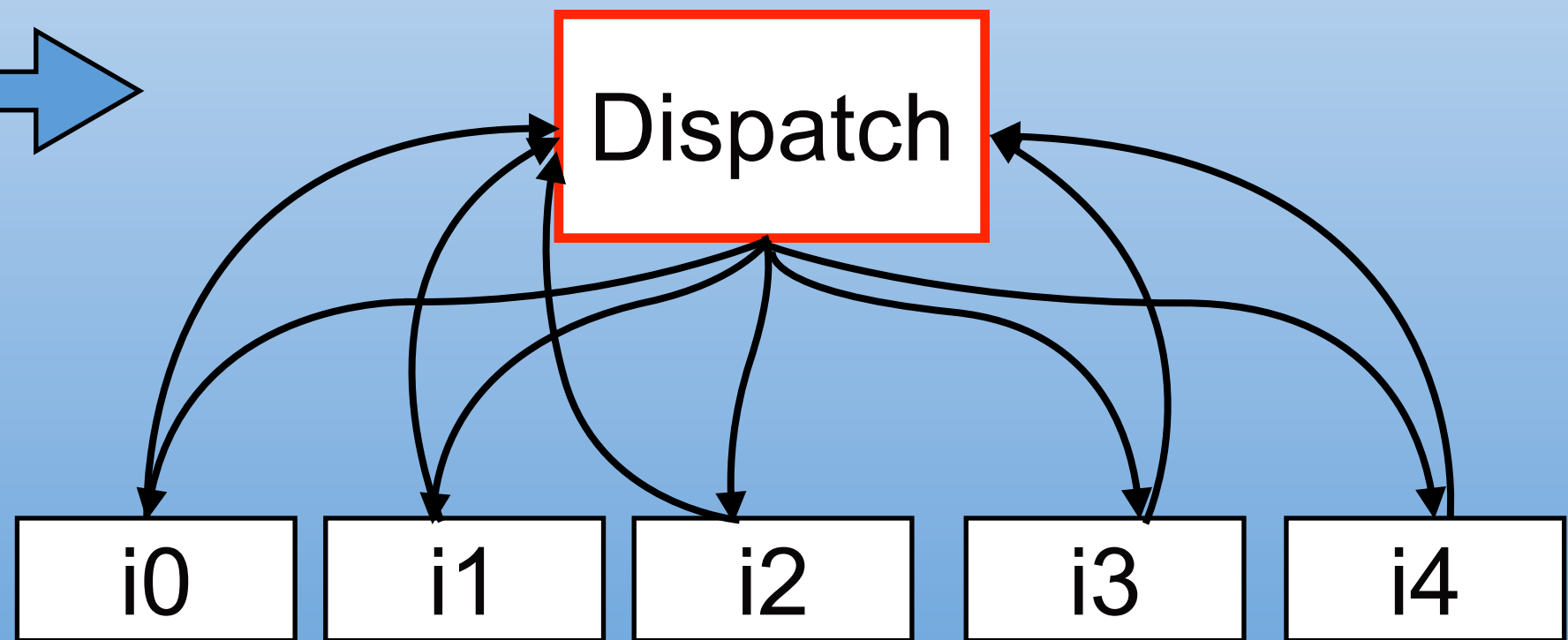
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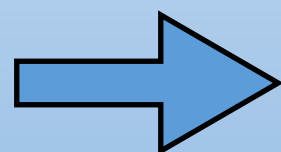
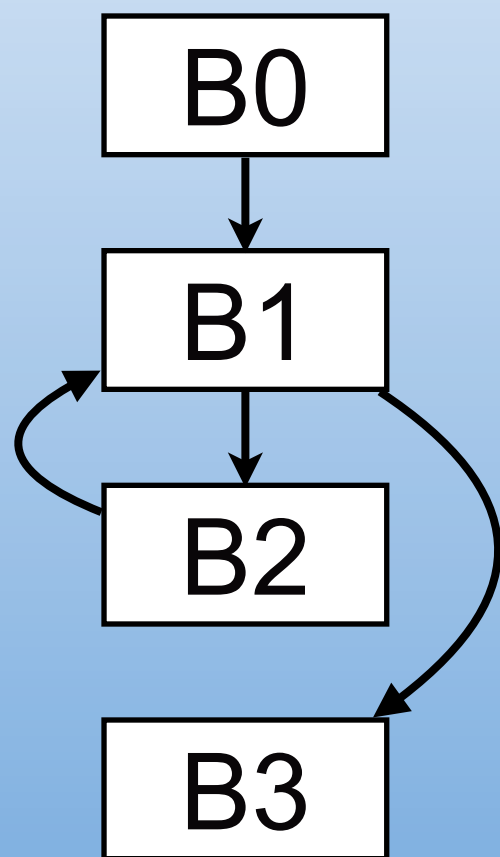
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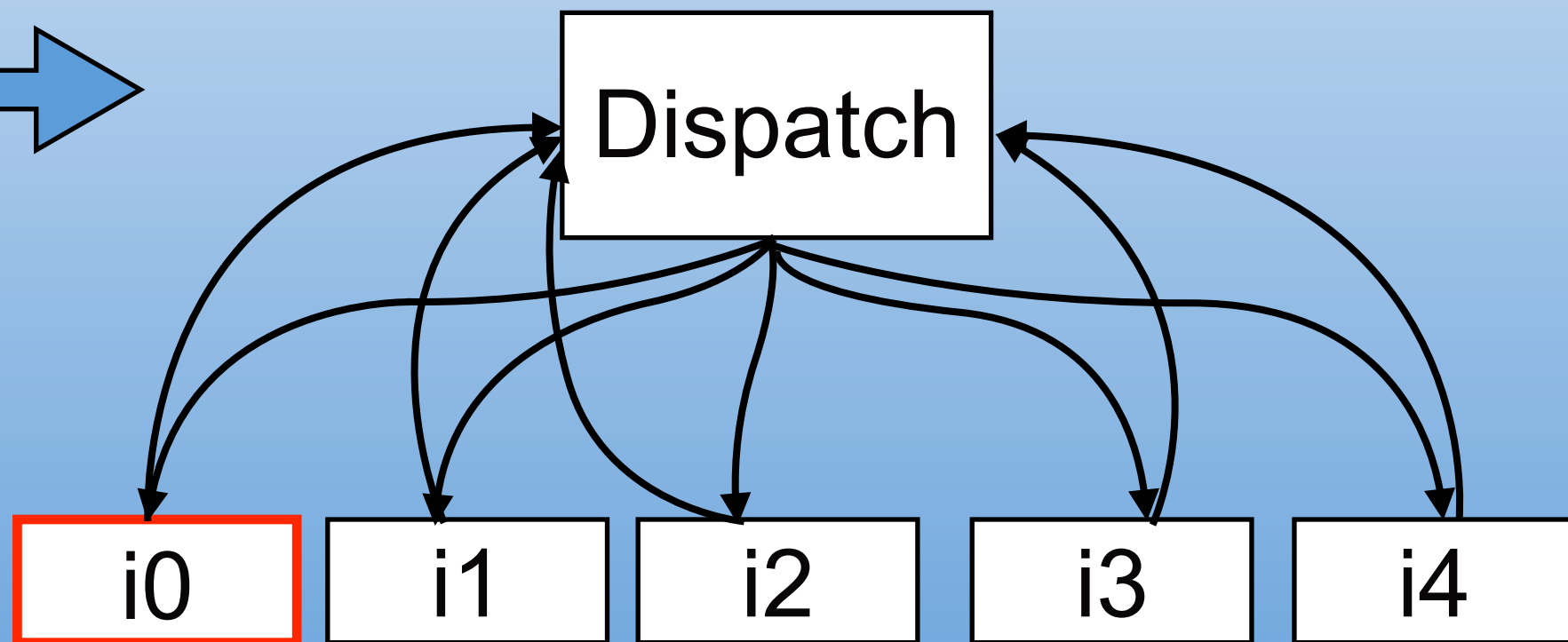
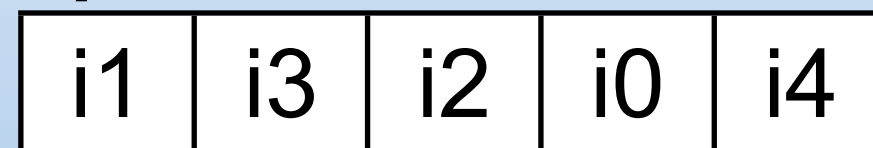
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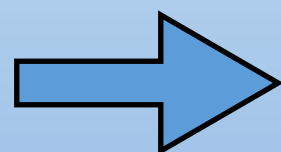
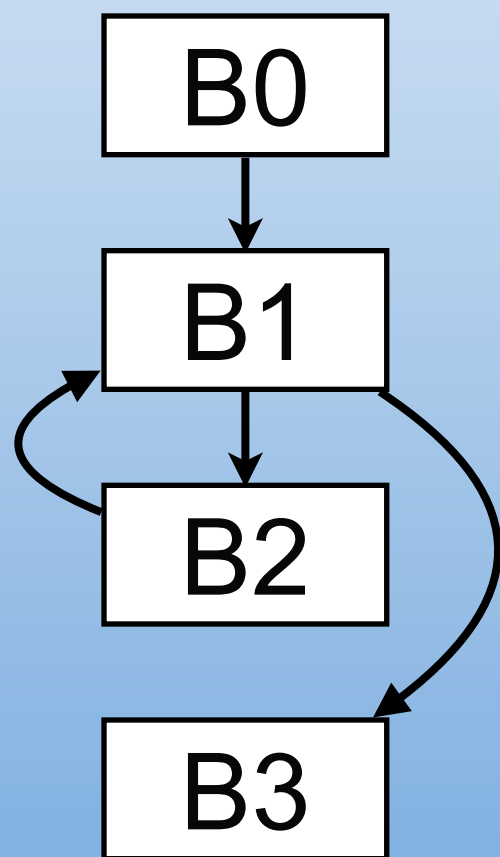
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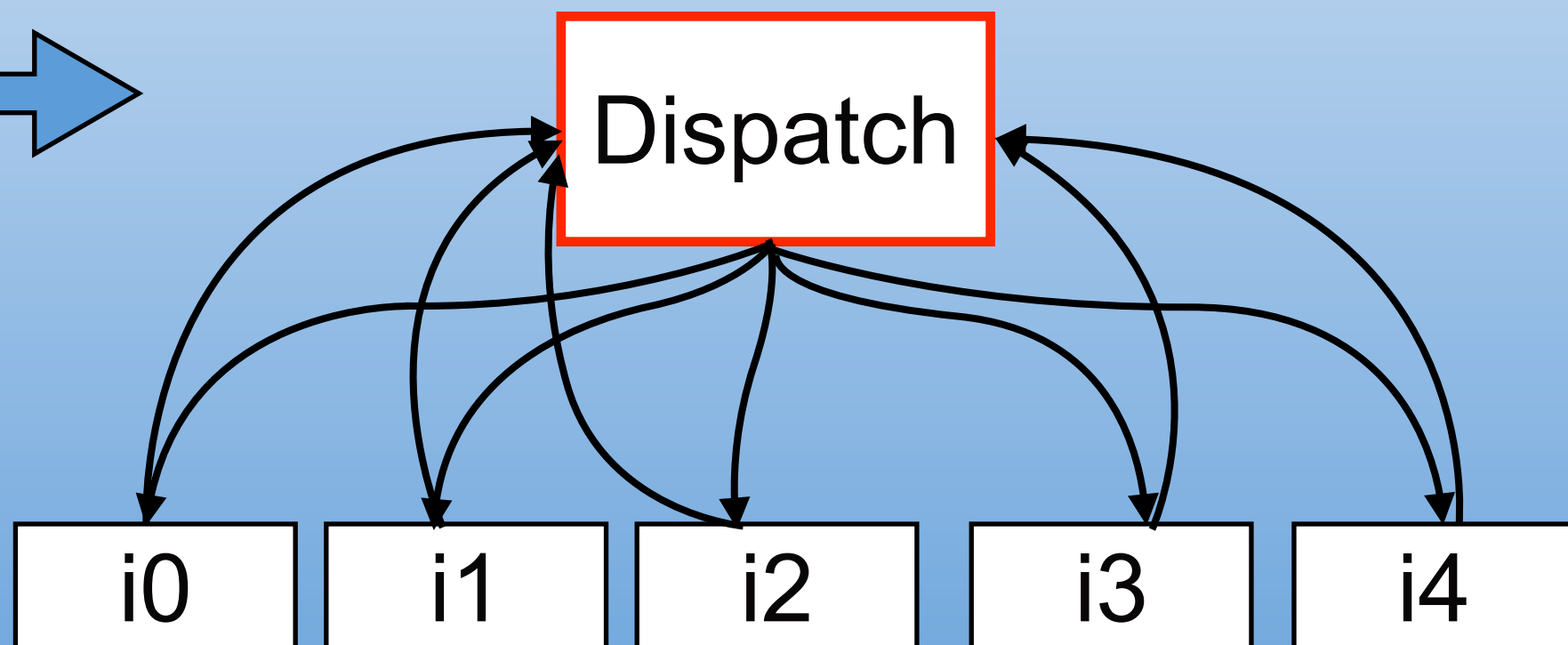
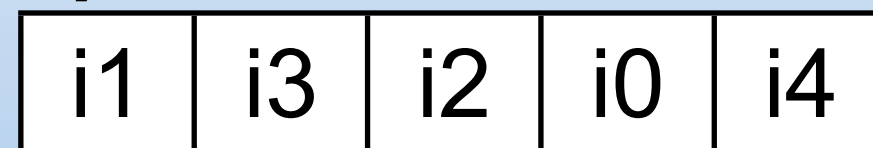
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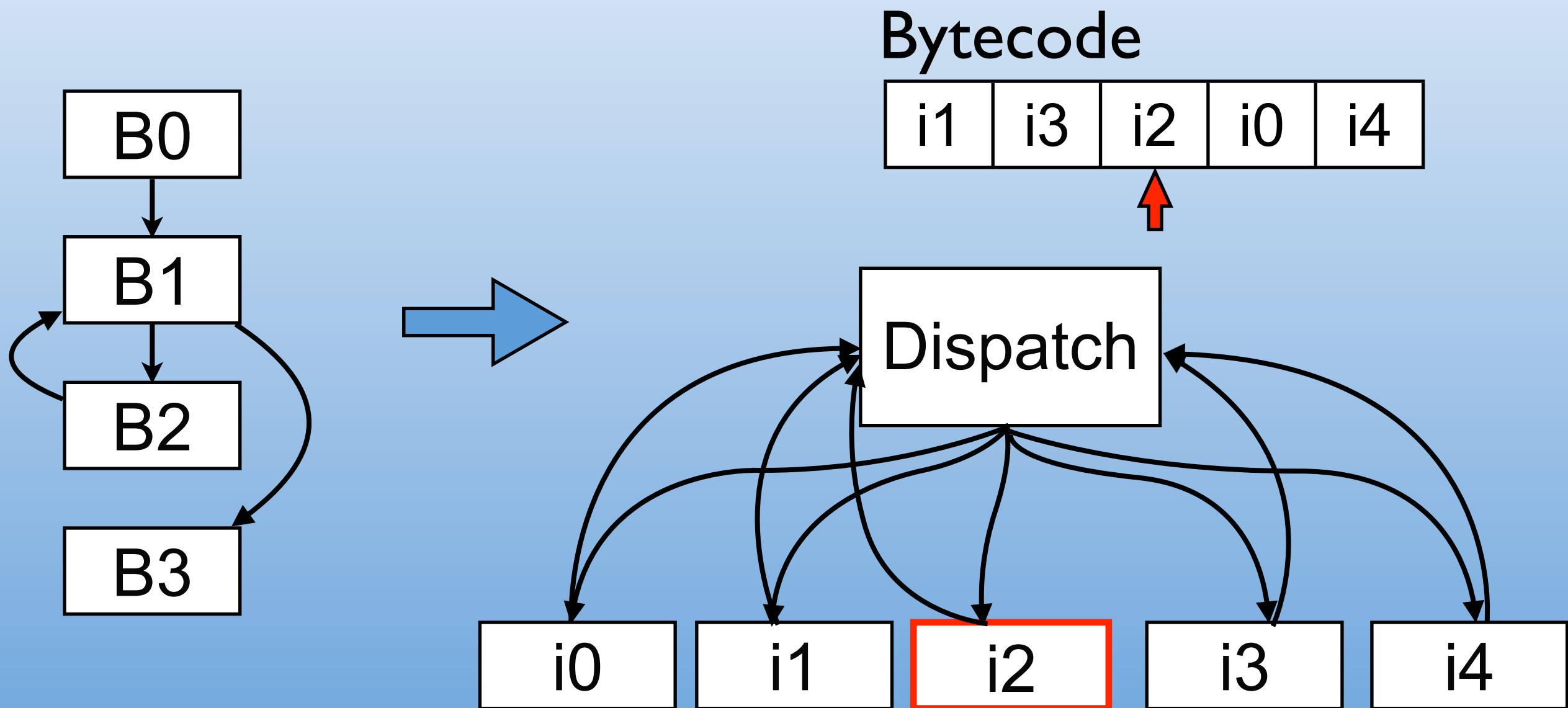
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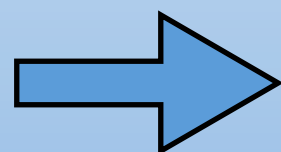
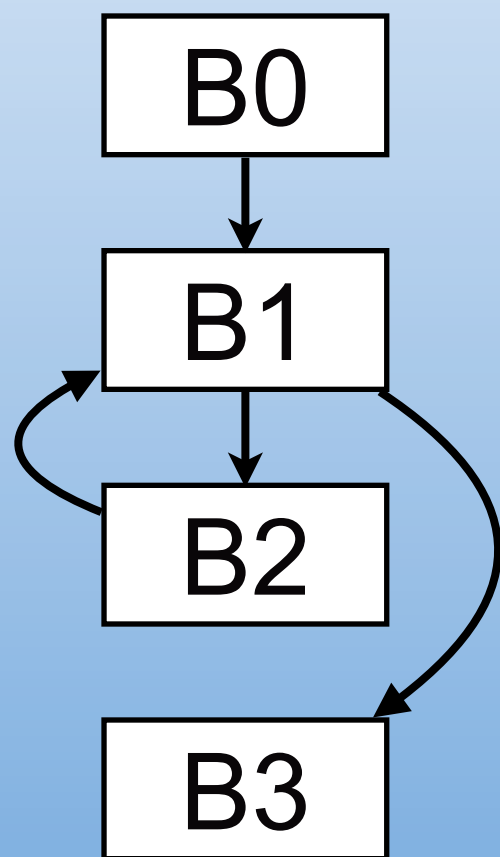
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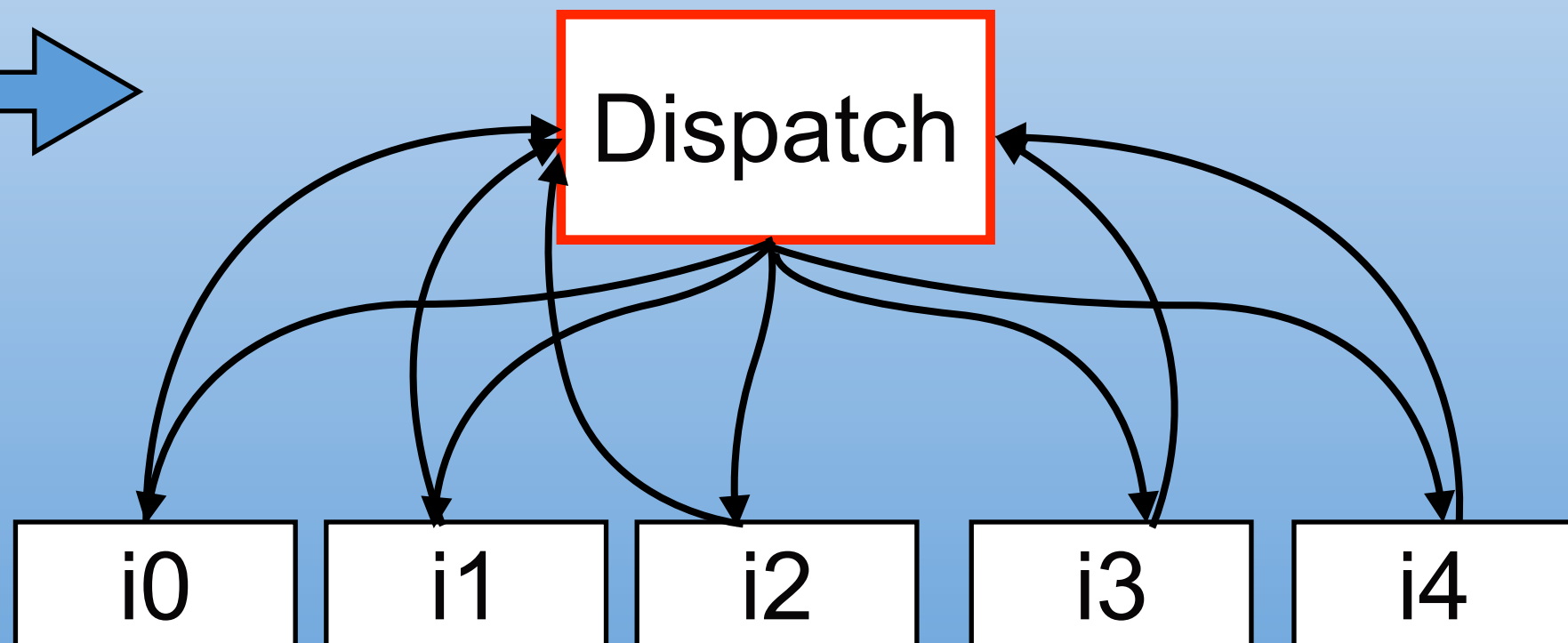
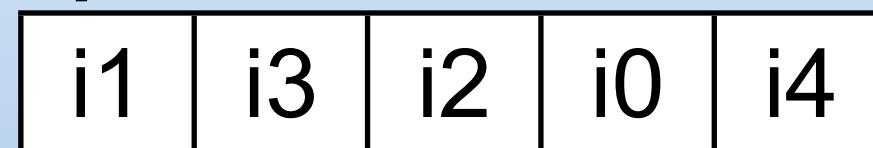
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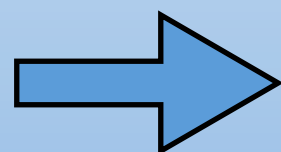
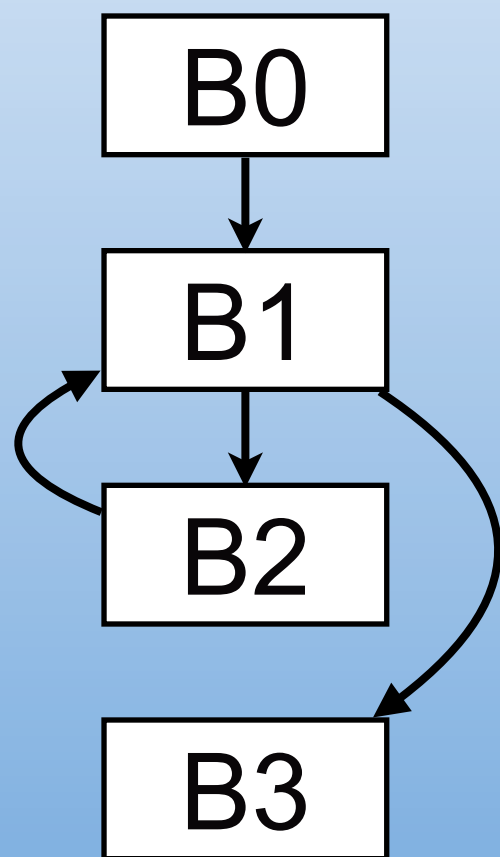
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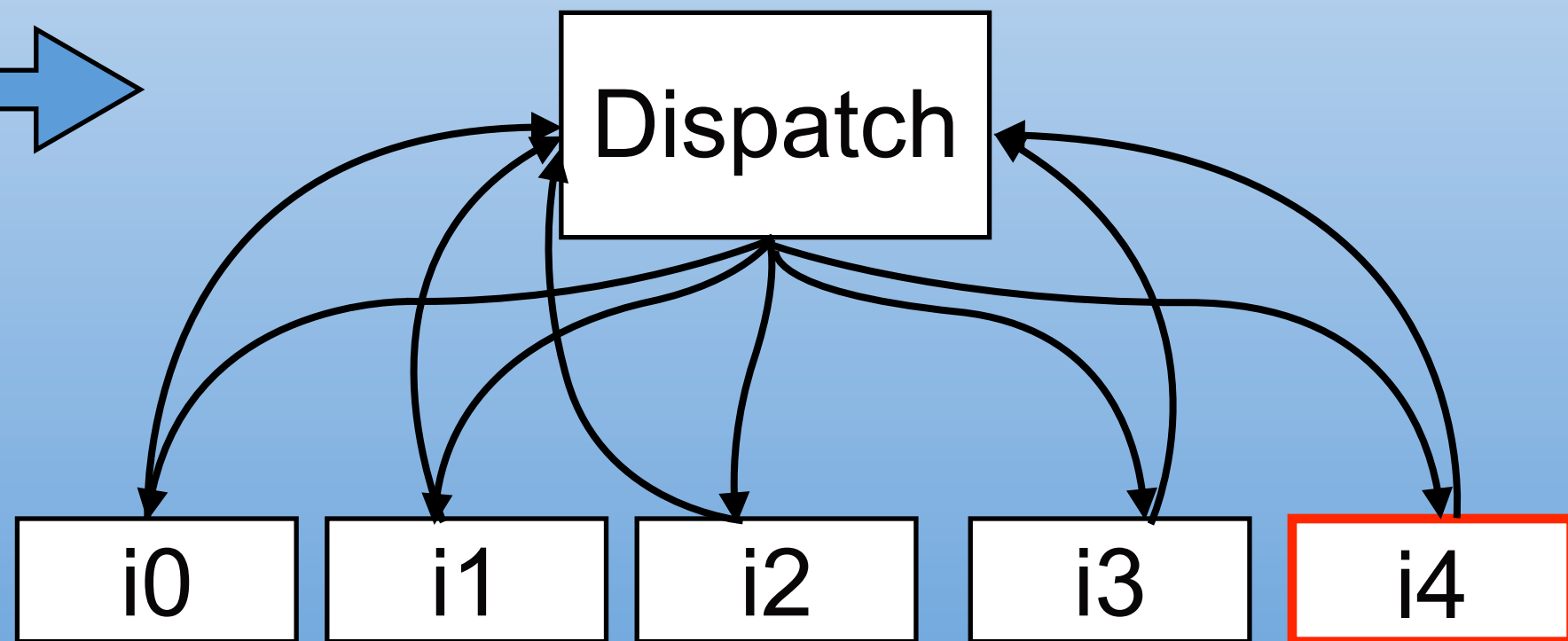
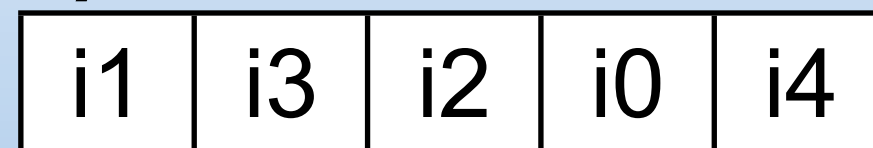
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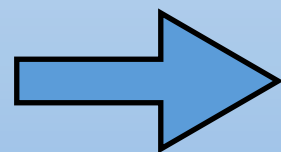
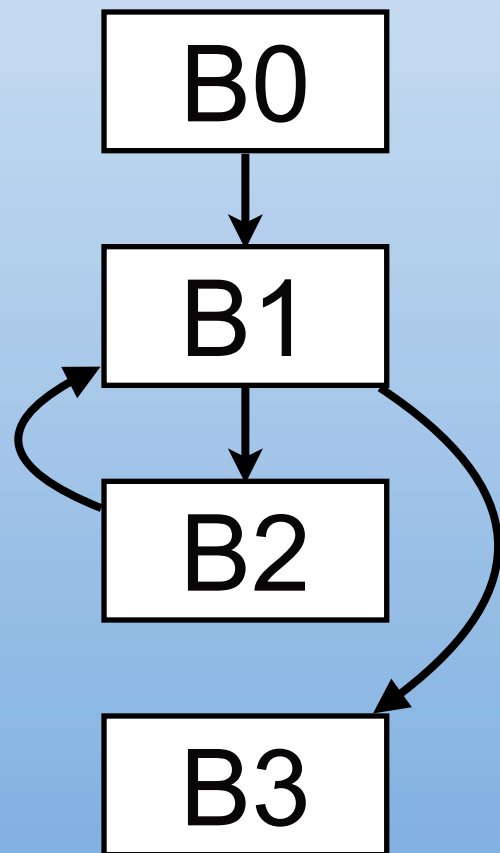
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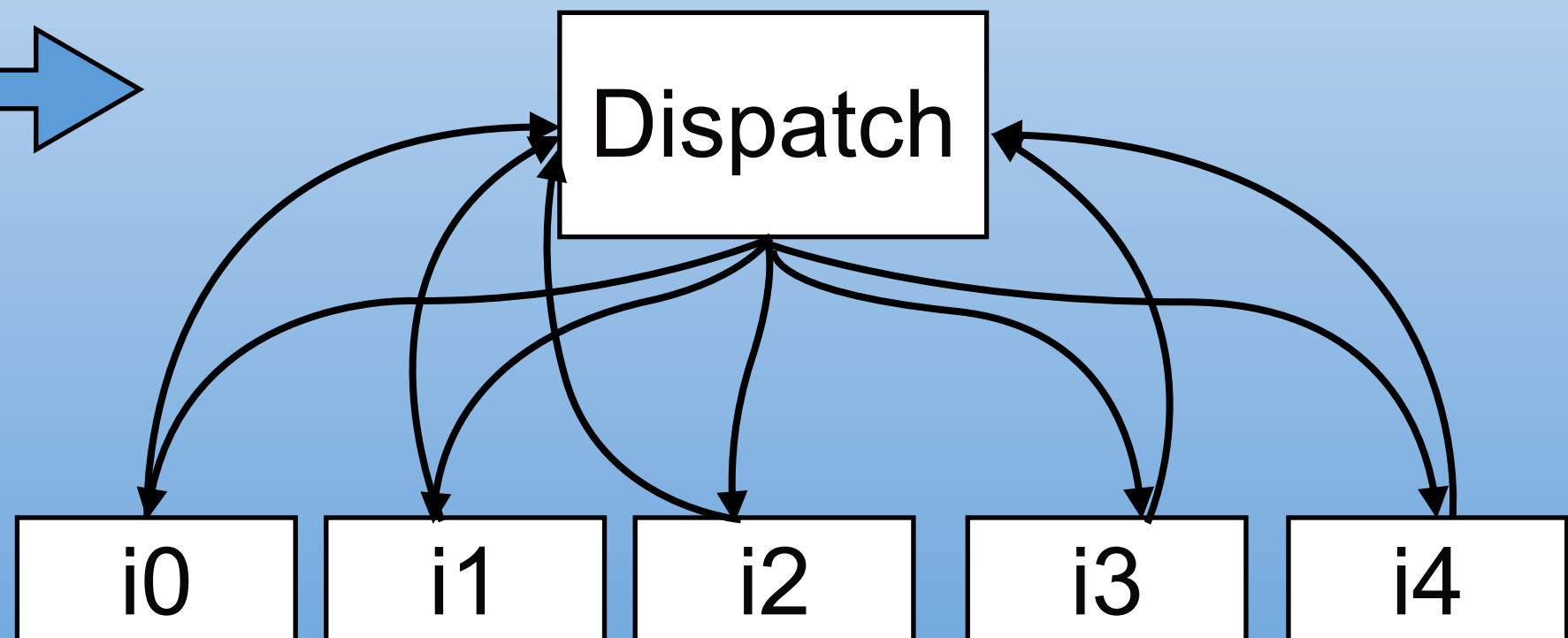
Bytecode



Emulation-Based Obfuscation



Bytecode



Emulation-Based Obfuscation

- Fundamentally different from string obfuscation
- Original program is hidden under interpreter's logic
- Looks like normal, unobfuscated program

Anti-analysis Defense

- Whether to conduct attack is determined at runtime
- Detecting detection/analysis systems:
 - *behavioral differences*
 - *execution overhead*
 - ...

Implicit Conditionals

- Hiding possible branch targets from multi-path exploration
 - branch without conditionals
- Emulation-based obfuscation?

Implicit Conditionals

- Replacing conditionals with calculation of parameters used by the interpreter
- ~~Emulation-based~~ *conditional-free* obfuscation
 - *semantics-preserving*

Implicit Conditionals

$$\begin{array}{l} \text{if } \text{test}(v) \\ \text{then } C \end{array} \xrightarrow{\text{transform}} \begin{array}{l} p = \text{comp}(v) \\ \text{interpret}(p, \text{code}) \end{array}$$

- $\text{comp}(v)$ returns desired value *iff* $\text{test}(v)$ holds true
- code is the bytecode representation of C , which is arranged based on the desired value of p

Possible choices of parameter:

- Bytecode entry point
- Instruction pointer increment value

Implicit Conditionals

Example:

```
bytecode = [b1, b2, ...];  
p = hash(navigator.userAgent);  
ip = parseInt(p[4] + p[11]);  
inc = parseInt(p[13] + p[7]);  
while (ip < bytecode.length) {  
    execute(bytecode[ip]);  
    ip = ip + inc;  
}
```

Experimental Evaluation

▶ Target systems

- VirusTotal

www.virustotal.com

- Zozzle

Curtsinger et al.

- Wepawet

Cova et al.

▶ Testcases

- 7 malware samples

- 2 versions for each sample

- code unpacking
- proposed protection

Detection Results

malware Sample	Code Unpacking			Proposed Protection		
	VirusTotal	Wepawet	Zozzle	VirusTotal	Wepawet	Zozzle
M_1	5/40	detected	not detected	0/42	not detected	not detected
M_2	4/41	detected	not detected	0/42	not detected	not detected
M_3	5/42	detected	detected	0/42	not detected	not detected
M_4	5/42	detected	detected	0/42	not detected	not detected
M_5	5/42	detected	not detected	0/42	not detected	not detected
M_6	5/42	detected	not detected	0/42	not detected	not detected
M_7	10/42	detected	n/a	2/42	not detected	n/a

Discussion

- Can we adjust existing tools against proposed protection?
- Possible attack models

Conclusion

- Existing detection techniques are closely tied to current malware and protection.
- Can we defend future attacks regardless of obfuscation/protection?

Thank you